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# DESIGN AND IMPLEMENTATION OF HIGH SPEED AES ALGORITHM FOR DATA SECURITY

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# ABSTRACT

In the current cyber era, 44% of the Indians use internet today and the numbers are increasing exponentially in recent years. It is accepted truth that in the coming year's security architecture in cyberspace decides growth of nation.

In this paper, we present a VLSI based AES (Advanced Encryption Standard) encryption that effectively addresses espionage and fraudulent cybercrime based cyber attacks. It is most commonly used symmetric block cipher algorithm that transform information into obscure data based on key-defined transformation set. In addition, it is lossless operation with size of input and output being the same and could be extended to a wide range of applications. We limit our focus to128-bit AES encoding and decoding operations over VHDL coded transformation that requires key for successful completion of the operation With rise in several methods proposed for implementation of data security it has become more vital for a feasibility study of any hardware design is essential i.e. to test encryption and decryption process of the proposed 128-bit AES algorithm. In the simulations results, we analyze the each of the transformation that is incorporated for coding on FPGA using Xilinx ISE tool.

**KEYWORDS**: Encryption, FPGA, Advance Encryption Standard, Xilinx, Data Security.

### INTRODUCTION

"Digital India" program significant portion of government records, personnel information, and financial details of Indians would available online [1]. Its main aim was to enhance the electronic literacy among the people so that they could acknowledge the concern government benefits which in turn reduce corruption. Unfortunately, every person's financial, medical, social and criminal history would be online and digitally available for individuals of both nature (i.e. good and bad). Henceforth, the demand for the data security has been growing exponentially now-a-days. Fundamentally data security deals with the two practices for developing algorithms namely:

- 1. For preserving the integrity of data in consideration
- 2. For preserving the information of data in consideration

Even though both the practices sound similar there is significant difference between both i.e. one preserves the integrity of the information by detecting any changes in data (for example: hashing algorithms). Most of the world is dependent upon software that protects the information of data (for example: encryption). This fact signifies the predominance of security activities in the current digital and cyber era. In general these algorithms will be developed with series of different activities which need to be carried out simultaneously [2]. With the recent developments in the digital technology have increased the ease with which multimedia data are reproduced and retransmitted. These developments have increased the potential of data manipulation and alteration of digital data with little distortion. Typically, a covert communication system can be viewed as secured communication system wherein the data that is being communicated in unintelligible nature. Cryptography is one of the approaches that provide the solution for a covert communication system based on mathematical and/or statistical transformation. Existing algorithms for cryptography can be broadly classified into two classes depending on cipher code generated:

1) Stream cipher and

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2) Block cipher

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Most of the researchers have focused on transformation in the block cipher domain as it attains a higher robustness. Block cipher based cryptography is one of the active research area, that develop methods to protect against privacy of multimedia and visual surveillance system for security proposes. Henceforth, we a brief review on each of systems in a simplified manner with basic encryption and decryption process is illustrated in figure 1.



Figure 1. The basic encryption and decryption process

From the above, *plaintext* refers to original critical data that needs to secured against espionage and fraudulent cyber attack. The obscure data called *ciphertext* is obtained after a series of transformation known to be a part of the encryption process. On other hand, the decryption performs the series of inverse transformation over the *ciphertext* to obtain plaintext without any distortion. The series of transformation employed to secure the digital information often encompass mathematics, applied statistics and code programming.

It is evident fact that there is no clear definition on the performance of the encryption that could be calculated. Therefore, the performance analysis without some logically accurate approach is employed for analysis based on set of well-defined and bench marked parameters. In addition, a perfect coded cryptosystem is not only flexible in the security mechanism, but also has high overall performance based on processing time, delay and size (area). In spite of several short comings, the hardware based encryption approaches have evolved drastically and gave precise security to the communication network. It is a documented fact that if the problem domain and corresponding encryption factors (key) are not well defined then a proper encryption of the entire data is not possible. Henceforth, there will be unforeseen contingencies that would be considered due to the miss-communicated, non-communicated, short or vague description of hardware requirements. Several issues that plague the progress of the effort are the common understanding which it very complex and highly improbable to precisely assess the entire design to development. The primary steps in any design process were to comprehend, evaluate, analyze and characterize the system based on the requirement and application in which it would be incorporated. When proposed architecture of the hardware is in sufficient or when the performance of hardware is limited then more often we alter the associated coding (i.e. software) in design to ensure proper working of the system as expected.

It is quite evident that the numbers of resources for exploiting an encryption within the constraints of hardware are limited; all features as per the requirement are generally not included due various constraints. In addition, any encryption technique comes with considerable risks that can be mitigated as the most vital and critical attributes such as the key management. It is essential for us to optimize the resource allocation which has an important effect on effectiveness of proposed design architecture. Further in a typical & practical aspect, for any organization to assign better and optimized resources for complex and vital information. Any encryption system is generally evaluated throughout the entire stages of operation i.e. encoding, transmission, receiving and decoding.

In brief, the proposed methodology focuses on implementation of adaptive and robust 128-bit AES encryption algorithm in hardware architecture. The proposed system offers highest level of security for various applications such as access control, secured communication and etc. It is essential to comprehend the holistic capability of the organization factors to offer the expected solution within specified resources and delay. Therefore, there must be a meticulous procedure that permits a detail analysis and considers other parameters within the problem domain.



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In this section, we consolidate the brief background on various attributes and functionalities that are factored in designing the architectures of existing AES encryption technique while addressing hardware issues such as area, power and throughput. Although some these are very popular than others techniques that address data security issues, the constraints associated with performance, size and power that differentiate each of these architectures. It is evident that each technique has several advantages and corresponding limitations, so to obtain an precise analysis all techniques needs to be employed that will make the entire process complex and very costly. The excellence and scope of cryptography based approaches doesn't limit to the securing the digital assets in consideration, but based on the time taken for the breaking the obscure information into meaning data by malicious eavesdropper. AES (Advanced Encryption Standard) is the popular eminent algorithmic frameworks that are commonly employed for effective encryption whose key parameters can be designed that are adaptive to application based development platform, which is vital aspect for encryption algorithm [4].

Advanced Encryption Standard (AES) (commonly coined as Rijndael) is a symmetric block cipher standardized by NIST for securing information. In this thesis, we limit to design and implementation of 128-bit AES encoding and decoding techniques wherein VHDL coding is employed for defining the transformations during encoding phase and inverse transformation during the decoding phase of the system. The synthesized analysis was simulated Virtex-2 based FPGA while feasibility of the approach is studied using the Xilinx. The transformations are defined as the set of executable code which performs a computational complex operations in an optimum manner as per the application requirement. It is vital that the outcome of the system needs to be consistent, efficient and reliable based on the key applied within the application in consideration [14]. The requirement of the data security is evolving into an active research area due to the rapid changes in processor speed and application requirements in consideration that operates the entire system.

Gnanambika, M., Adilakshmi, S., and Noorbasha, D. F [5] proposed hardware implementation of AES that offers high speed and efficiently addresses the Sub Byte transformation method. This exploits the benefits of combinational logic performed in polynomial bases while the architecture was implemented by combining composite field arithmetic in normal bases. Furthermore, key expansion was also presented with a detail and efficient architecture simulated using Modelsim proves the efficient performance of the system. Kalaiselvi, K., and Mangalam, H. [6] proposed AES technique with a key expansion process to address low power and high throughput requirements of hardware design. The overall power consumption of the encryption system is reduced while the delay along the critical path is also reduced due to optimized architecture. Simulation analysis proves that this architecture offered superior performance in comparison with other existing architectures in terms of power, throughput and delay. In addition, 256-bit key is employed during encoding and decoding process.

Das, S. S., and Resmi, R. [7] introduced Advanced Encryption Standard (AES) based VLSI implementation with an effective and optimized architecture with a module that exploits the key expansion with ROM. The incorporation of ROM at the cost of commonly used is quite beneficial. Furthermore in an effort to reduce the area and power of the system registers a novel framework of merging two steps with exclusion of shift rows was studied. The Virtex5 FPGA was utilized during simulation analysis to study the feasibility modified AES algorithm. The results prove that they achieve a higher FPGA efficiency (in terms of Throughput and Area) compared to previous existing AES designs.

Dattathreya, K. A., and Kashwan, K. R. [8] developed AES with a multimode multiplier capable to compute polynomial fields. The Montgomery Multiplication are complex and operational time is high which makes AES algorithm slow, and power consumed is high that was effectively addressed using a Wallace Tree multiplier. The delays along the critical paths of the design are optimized by exploiting buffer insertion concept. In addition, design analysis and synthesis were simulated using Xilinx while the coding was done on Hardware Description Language (HDL). The results prove that they achieve desired efficiency in comparison with previous existing AES designs and model.

Meera, R., and Kalaiselvi, K. [9] proposed a 128-bit AES implementation on FPGA that could address information security issues that has become a prominent area now-a-days. They illustrate various advantages and limitations associated with implementation with aspect to hardware and software. The results prove that they 128-bit AES



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approach offers a considerable high efficiency and reduces the computational time taken for key generation. In addition, the system achieves desired speed while the low area and power are maintained.

Vaidehi, M., and Rabi, B. J. [10] proposed a novel implementation of AES with modified architecture that reduces the complexity of MixColumn transformation along the critical path. In addition, AES architecture is optimized by innovative re-designing of MixColumn transformation by minimizing the repetitive logical functions. Thus, this system addressed the delay, power in an effective manner than the traditional methods realized. VHDL is exploited in design phase of the system but the Common Sub-expression Elimination (CSE) incorporated to significant and effectively optimizes the MixColumn transformation of algorithm. Further, design improvement stands at 10.93% (hardware slices), 13.6% (LUTs) and finally 1.19% (delay). The simulation analysis shows some of the promising results with 4.75% (area), 4.56% (power) in comparison with traditional MixColumn based AES Encryption. Zhang, Q. et.al [11] proposed a novel implementation of asynchronous AES with modified architecture with a round key generation reduces the complexity and delay along the critical path. Even though it is most prominent encryption technique used as security for financial applications, unfortunately it is not immune to side-channel attacks which are addressed by the Asynchronous AES design. In brief, this system incorporates a combination of round key generation and mix column calculation to implement novel design. Furthermore, to reduce the area and power of the system Balsa HDL properties are incorporated on the basis of GTECH-based design flow. Experimental results prove that this system offers better performance i.e. 67.7% (lower computation time) and 40% (lower power) while the 7.3% (area) and 15% (delay) are reported.

Kshirsagar, R. V., and Vyawahare, M. V [12] proposed AES architecture by partitioning of buffer into sub-blocks of pipeline structure with repetitive AES modules. They are utilized intermediate buffers along with the module of byte substitute operated over shift rows block (Shift Row). This design improves the area with optimized usage of blocks while the integrity of the AES encryption algorithm is preserved. The feasibility of the proposed system was carried out on Xilinx's SPARTAN-3 FPGA that shows promising performance in terms of throughput rate and hardware area. The simulation results prove that this system offers 56% (area) and 4.25% (delay) in comparison with the existing algorithms.

Due to the constant exploration within the field of Hardware engineering, there are several commonly used approaches were proposed for various classes of AES encryption algorithms. These methods available are classified into analogy based various factors. Based on the simulation results, we observe that each of existing approach has advantages as well as drawback in comparison with other approaches, as their advantages and limitations are often complimentary to each other [13-15]. To understand the advantages and limitation of any method it is very important to know when you can use which method to optimize the resources while maintaining the performance. It is an evident fact that hardware design is complex problem that doesn't has a precise solution that could be easily deciphered because there are several parameters used in the computation for a precise solution i.e. AES encryption algorithm. The prior understanding helps us in optimizing the resource allocation which has an important effect on effective implementation of the AES architecture. Regular revaluation is a proven manner to ensure the advancement in design and supervise the system performance.

#### ADVANCED ENCRYPTION STANDARD ALGORITHM

Any encryption method consists of two vital sub-blocks that decides its performance and strength i.e., the algorithm and secret key. As discussed in prior sections, algorithms are usually function based on complex mathematics and/or statistics with the scope of defined constraints. Whereas the key is a collection of bits, that are essential at various stages of the algorithm for decoding the information from obscure data. Most of the existing encryption methods can be classified based on the characteristics of algorithm and key. The entire integrity of the system depends on the protection of the key as it compromises both the encoding and decoding process as presented in figure 2.



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Figure 2.Schematic of AES

Advanced Encryption Standard (AES) is a popular eminent symmetric block cipher that converts electronic data to an obscure form which protect integrity of data from eavesdropper whose key parameters can be designed that are adaptive to application based development platform, which is vital aspect for encryption algorithm. The prime issues that are addressed by it are namely: security, efficient and exportable. In this thesis, we limit our focus to 128-bit AES encoding and decoding block and process is symmetric in nature i.e. same binary key is employed during both processes. It is a simpler lossless reversible algorithm that operates on bytes. The AES comprises of different transformations that operate in an iterated block cipher over the 128-bit fixed block size and a variable length secret key. Each stage is array of which are processed based on the sequence of transformations dictated by the secret key of varying length. AES utilizes round function for both encoding and decoding based of four different transformations:

- 1) Byte substitution using a substitution table (SubBytes)
- 2) Shifting rows of the State array by different offsets (ShiftRows)
- 3) Mixing the data within each column of the State array (MixColumns)
- 4) Adding a Round Key to the State (AddRoundKey)

### **ENCRYPTION**

In this section, we present the basic procedure for encoding data based on AES algorithm. This algorithm generates a binary blocks of 128-bits using the specified input parameters, which are XOR with input key to generate the input for round-1 as the number of rounds increases transformations applied based on various logical and/or non-logical operations becomes complex. In addition a parameterized round function is exploited for key schedule that consists of a one-dimensional array of four-byte words derived using the Key Expansion routine. The individual transformations that carried out are namely SubBytes, ShiftRows, MixColumns, AddRoundKey.

**SubBytes Transformation:** The SubBytes transformation is a reversible, affine based non-linear byte substitution with each byte transformed independently. Figure 3 illustrates the effect of the **SubBytes** transformation in hexadecimal form in detail.



Figure 3. SubBytes Operation of the State



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**ShiftRows Transformation**: This transformation shifts each row in a circular manner rather than the bit wise. The circular shift ensures each position is altered i.e. the shift of last position bit would be the first position in the same row, depending on the row index which is illustrated in the figure 4.



Figure 4. ShiftRows transformation

**MixColumns Transformation:** In MixColumns transformation focuses on column-wise which are multiplied against the matrix. The results of these multiplications are XORed together to produce inputs for the next stage. The predefined 4X4 matrix value and the first column of the ShiftRows state are represented as follows, for the multiplication.

S <sub>0,c</sub>		02	03	01	01	S <sub>0,c</sub>
s'1,0		01	02	03	01	<i>S</i> <sub>1,c</sub>
s' <sub>2,c</sub>	-	01	01	02	03	\$2,c
s' <sub>3,c</sub>		03	01	01	02	S3,c

**AddRoundKey Transformation:** The **AddRoundKey** transformation performs a bitwise XOR operation between the Round Key and the data block. In addition, we need to ensure that the bytes obtained during the expansion of key should not be used. For example, let's assume that the initial 16 bytes are XORed with expanded key then those bits are never employed in any other transformation.



Figure 5. AddRoundKey transformation

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# [Mishra\* et al., 5(8): August, 2016] IC<sup>TM</sup> Value: 3.00 KEY EXPANSION

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It is general understanding that the key must be expanded prior to encoding and/or decoding process. Addround Key transformation discussed in earlier section supports the key expansion that is performed at the initial stage of the coding process. In general, the expansion of key ensures the size of the key is more than the sub-keys that often provides enough space for the algorithm. The key size and ExpandedKey size are interrelated and largely depend on the number of rounds. Most of the key expansions are performed based on few major routines that are executed in consecutive that are namely

- ROT WORD (4 bytes: circular shift)
- SUB WORD (4bytes: substitution)
- > RCON
- > XOR

#### DECRYPTION

The decoding process is a straightforward process where all necessary parameters are dictated based on the secret key. The Cipher transformations are simple inverse application of the encoding process at each for the AES algorithm that are listed below

- InvShiftRows
- > InvSubBytes
- InvMixColumns
- AddRoundKey

# **PROPOSED ARCHITECTURE**

In recent years, various designs were introduced for AES of different key combinations to address security based applications; unfortunately their robustness and immunity against brute-force are not analyzed but the delays and power consumption are given priority. The major problem regarding the AES lies with complex transformation that includes repetitive transformation and shift operations that induce significant delay and require high computation time to complete the operation thus consuming high power. In terms of hardware engineering, the measure of hardware design quality and feasibility can be expressed as a function of the issues related with the product i.e. is it has minimum delay, size, and power consumption and so on.

With significant rise in the demand of the quality in processor, AES based architecture is becoming a popular method to address such issues of data security. In the last few years, the main motto of the researchers was to enhance speed of the AES operation in consideration. Each block has several advantages and corresponding limitations, so to obtain a precise execution at end of each stage we need to focus on the performance of each module. The marking of the individual module/block based design to investigate various Pipelined transformations to achieve low power consumption while offering high speed which is prime motivation of this investigation. The feasibility study of the entire AES modeled on ModelSim depends on three factors i.e. size, power and delay which are calculated for a transformation stage, thus it helps us in designing high speed AES architecture in an optimized and effective manner. The state is defined by the block's current condition and the transformation alters the block at each round.

#### **Implementation Requirements**

The proposed system is designed based the advantages on various existing architectures and the quality set to address complicated problems. As per the outcomes, the constraints are projected with detail implementation requirement was used for alternative creating and subjective judgments on hardware design. Implementation requirement would enhance alternative creating and consolidate it with security activity and computational zing soon think about the interrelations between delays, complexity metrics with repetitiveness, and power consumption.

**Data Length Requirements:** Any data that could be converted into a bit stream can act as input data for AES algorithm based implementation but we length is fixed (i.e.128bits). The data length is fixed for both Encoding and decoding blocks of AES implementation i.e. 128 bits.

*Key Length Requirements:* In this paper, the key length is fixed to 128-bits but there various key lengths that could be incorporated.



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*Keying Restrictions:* There are no constraints or restrictions in the selection of the secret key but care has to been taken while communicating the key among the authorized users.

*Parameterization:* key length = 128 bits, Nb (size of block) = 4 and Nr = 10.

#### Implementation Flow AnalysiS

Based on review of the prior market research, shows that most of the existing designs and their requirements are focused on capitalizing the repetitive features of the algorithm. Based on this detailed survey, the specific feasibility of the encoding and decoding architectures was identified for the RTL modeling. The FPGA offers inter connected logic blocks with a two dimensional arrays wherein both are programmable. Transformation block are programmed to implement a desired function and interconnect are programmed using the switch boxes to connect the transformation blocks. Our basic metric set of FPGA focuses on size (area), computation time, power dissipation and unforeseen defects for the feasibility of the system, and optimized design.

The fundamental design based on the VHDL, schematic or XOR expression. The optimization of these expressions was carried out by considering area or speed of the system. These impacts contribute by all odds to the "dynamic engagement" property of compelling learning. It is evident that VHDL is synthesized and is directly proportional to estimated size, duration, environment in with it is executed, and quality of power source available to execute based on following steps

- > Translate
- > Map
- Place and Route



#### **Synthesis Result**

The alternative analysis is created RTL methodology, wherever it imparts their seeing concerning endeavor house. The RTL model is employed as basis in implementation of synthesis process on FPGA Virtex-2 board. The developed AES Encryption and Decryption Algorithm preponderantly targeted on functionality based mostly alternative creating by taking all the thoughts of the teams and structure the quality set. With proposed algorithm getting to be a lot of dynamic, self-ruling we evaluate the flow analysis of the proposed AES implementation would performed on Virtex-2 board along with the Xilinx tool and other concerns are addressed. In this phase, gate level based net-list is generated based on the proposed system RTL model that is mapped with SPL. The proposed high speed design of AES algorithm would be implemented on a Virtex-2 family.

#### **RTL** Schematic

The RTL (Register Transfer Logic) shows the inputs and outputs that can be viewed after synthesize of design as black box. The above figure 7 shows the top level block of the Encoding and decoding schematic that contains the primary inputs and outputs of the design.



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Figure 7. RTL Schematic

#### Device utilization summary:

Device Utiliza	ation Sumn	nary	
Logic Utilization	Used	Available	Utilization
Number of Slice Flip Flops	1,626	93,184	1%
Number of 4 input LUTs	161,979	93,184	173%
Logic Distribution			
Number of occupied Slices	81,991	46,592	175%
Number of Slices containing only related logic	66,234	81,991	80%
Number of Slices containing unrelated logic	15,757	81,991	19%
Total Number 4 input LUTs	162,232	93,184	174%
Number used as logic	161,979		
Number used as a route-thru	253		
Number of bonded IOBs	391	1,108	35%
IOB Flip Flops	381		
Number of GCLKs	1	16	6%
Total equivalent gate count for design	1,055,595		
Additional JTAG gate count for IOBs	18.768		

#### Timing Summary: Speed Grade: 5

- Minimum period: 52.716ns (Maximum Frequency: 18.970MHz)
- *Minimum input arrival time before clock: 20.103ns*
- Maximum output required time after clock: 4.840ns
- Maximum combinational path delay: No path found

In timing summary, details regarding time period and frequency is shown are approximate while synthesize. After place and routing is over, we get the exact timing summery. Hence the maximum operating frequency of this synthesized design is given as 18.970 MHz and the minimum period as 52.719 ns. OFFSET IN is the minimum input arrival time before clock and OFFSET OUT is maximum output required time after clock. The device utilization summary includes Logic utilization and distribution along with gate count of the architecture. The summary also shows the details of resources used and available in %. Hence as the result of the synthesis process, the device utilization in the used device and package is shown above.

#### SIMULATION ANALYSIS

In this section, we perform simulation analysis on design and implementation of 128-bit AES encoding and decoding process. Computer simulations were simulated using ModelSim for conceptual analysis and Xilinx for feasibility analysis of the concept. The transformations are defined as the set of executable code which performs a computational complex operations in an optimum manner as per the application requirement. In order to test the robustness of the



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proposed implementation of AES that needs to be modeled in various scenarios and definitions of inputs and key. It is vital that the outcome of the system needs to be consistent, efficient and reliable based on the key applied within the application in consideration. Henceforth, Modelsim tool is employed to test the feasibility of the approach with reference to the functionalities and architecture characteristics of the design. Furthermore, AES based encoding and decoding operations are performed on fixed key size of 128 bits so as to obtained cipher text that should obscure with reference to original data which validates that proposed design works in an effective manner and as per the constraints defined by the algorithm.

#### Analysis

Initial investigation was based on the test bench that automatically forces random inputs for the process in consideration. Hence, the proposed system robustness could be measured by evaluating the waveforms for the various cases associated with the design flow of proposed FPGA architecture. The AES implementation offers inter connected logic blocks with a two dimensional arrays that perform various complex transformations. The various cases of the simulation are discussed in detail.



Figure 8. Case-1: Simulation analysis of proposed AES algorithm architecture

*CASE-1:* In this case, the key size is 128-bits and the both encoding and decoding were performed with the basic and common inputs i.e. chip enable (ce), clock (clk), and reset (rst). It is evident from the figure 8, that the data\_in bits and data\_out bits of encryption are quite different hence AES encoding has been operate. Furthermore, the data\_in bits of encryption and data\_out bits of decryption are similar which proves that proposed implementation is reversible. *CASE-2:* This case presents the detail analysis of the internal operation performed by the encoding system of the AES. Each of the transformation with input and corresponding output of each round could be seen in the figure 9. The encoding process incorporates a series of transformations such as substitution, shifting and mixing of columns to obscure the information.



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Here / /test_aes_new_120/a0/alast/line_30/s_row	No Data-		003 (00 00 0					111E931	7D 85) (CB	BZ ZL 72} (	3D 2E 89 5F	1 (AF US U7 :	94))
Here have new 120/-0/dout_ist	No Data-	1100 00 00	003 (00 00 1	<u>10 00} {00 0</u>	100.00} {00	00 00 00}}		1839 25	84 TD}{U21	1C 03 FB} {	JU 11 85 97	119 6A UB :	32}}
/test_aes_new_128/a0/dout_valid	-No Data-							$\vdash$					
	Ma Data	000000000	2000000000	0000000000	0000			V0005041	D.00D.COOF	000110507	100400000		<u> </u>
Hest_aes_new_126/au/data_out	-No Data-	.000000000	1000000000	4000000000	0000			13925841	DUZULU9F	BUC118597	196AUB32		<u>+</u>
		1.1.1.1	1.1.1.1	1.1.1.1	1 I I I	1.1.1.1	1.1.1.1	1.1.1.1		1 I I I	1.1.1.1.	L I I I I	1.1

Figure 9. Internal Operation of encoding phase of proposed architecture

• Ret_acc.new, 128/a1/and_m_add         1             1332541022003F801785371543832             1             1332541022003F801785371543832             1             1332541022003F801785371543832             1             1332541022003F801785371543832             13502             13502541022003F801785371543832             13502             13502541022003F801785371543832             13502             13502541022003F801785371543832             13502             13502541022003F801785371543832             13502             13502541022003F801785371543832             13502             13502541022003F801745             13502574F3C             13502             13502574F3C             13502574F3C             13502574F3C             13502574F3C             13502574F3C             13502574F3C             13502574F3C             13502574F3C             11             11	Dec_Inputs														
• Aret_ace_new_128/a1/dw_vaid_u1         • OCCUPSEDCT15597155A0E32         • OccupseDCT15597         • OccupseDCT15597         • OccupseDCT15597         • OccupseDCT15597         • OccupseDCT1559         • OccupseDCT1559         • Occups	/test_aes_new_128/a1/enc_dec	0													
Bit - NetLass_new_128/a1/data_m         6504400         D3528410.02000FBPDC11653715640822           VetLass_new_128/a1/Ass         0001000         287F1515284ED225648F7158039CF4F3C           Decuption         0001000         287F1515284ED225648F7158039CF4F3C           Detwork_est_new_128/a1/key_in         0001000         287F1515284ED225648F7158039CF4F3C           Detwork_est_new_128/a1/key_in         0001000         287F1515284ED245648F7158039CF4F3C           Bit Ass_new_128/a1/key_in         0001000         287F1515284ED245648F7158039CF4F3C           Bit Ass_new_128/a1/key_in         0001000         287F1515284ED245648F7158039CF4F3C           Bit Ass_new_128/a1/Adu_vaid_00         (154 913 3)         (154 913 2)         (154 913 3)           Bit Ass_new_128/a1/Adu/mat_41/Loox         (160 00 Dito)	/test_aes_new_128/a1/din_valid	1													
♦         Nest_ass_new_128/a1/key         001002           B         Nest_ass_new_128/a1/key         0010002         287E1516286ED2A6A6F7158030CF4F3C         001000           Decuyption         0010002         287E1516286ED2A6A6F7158030CF4F3C         001000         287E1516286ED2A6A6F7158030CF4F3C           Decuyption         001000         287E1516286ED2A6A6F7158030CF4F3C         001000         287E1516286ED26A68F7158030CF4F3C           Hest_ass_new_128/a1/key_01         0011000         287E1516286ED25468F7158030CF4F3C         001000         001000           Round1         (113111         (10014F368)(CESE 258)(E13F0C 28)(E5 63 0 C 66))         001000         001000           Hest_ass_new_128/a1/A0ut/mc_41/x_10w         (176 55 100 00 00 00)(00 000)((E12 51 13 02)(58) A1 42 (E16 C 57 28) (H 64 03 9 E7))         001000           Hest_ass_new_128/a1/A01/ime_41/x_10w         (160 05 00 100 00 00)(00 000)((E1 26 13 02)(58) A1 42 (E16 C 57 28) (H 64 03 9 E7))         001000           Hest_ass_new_128/a1/A01/ime_41/x_10w         (160 147 78 E1 (100 00 00 00)(00 000)((E1 26 14 07 21) (196 C 53 28) (H 64 03 9 E7))         001000           Hest_ass_new_128/a1/A01/ime_41/x_10w         (147 78 E1 (100 00 00 00)(00 000)((E1 26 15 00 P2) (10 03 04 04) (97 E C 3 95))         0010000000000000000000000000000000000	· ⊞	69C4E0D	(3925	841D02D0	09FBDC1	8597196A	DB32								
Bit - Aret_aes_new_128/a1/key       0001020       287£1515284E02Ak6P7758809CF4F3C	🔶 /test_aes_new_128/a1/k_en	0													
Decuption         Decuption         Decuption         Decuption         Decuption           B → Atst_des_new_128/a1/ksy_0         (13 111 11 (10 014F9 A8) (C5 EE 25 89) (E1 3F 00 C8) (E6 63 0C A6))         (13 111 11 (10 014F9 A8) (C5 EE 25 89) (E1 3F 00 C8) (E6 63 0C A6))           B → Atst_des_new_128/a1/ksy_0         (13 111 11 (10 014F9 A8) (C5 EE 25 89) (E1 3F 00 C8) (E6 63 0C A6))         (13 111 11 (10 014F9 A8) (C5 EE 25 13) (C5 22 2C 72) (30 2E 95 F) (AF 09 07 94))           B → Atst_des_new_128/a1/ksy_01         (14 77 56 F3) (15 FAD C21) (28 012 22 41) (57 50 00 E1)         (14 77 36 F3) (15 FAD C21) (28 012 22 13 10 57 50 00 E1)           B → Atst_des_new_128/a1/a01/ine_41/x_row         (16 06 D 00 00 (00 00 00) (00 00 00 00) (00 00 00 00) (00 00 00) (00 00 00 0	· ⊕	00010203	2B7E15	1628AED2	A6ABF715	8809CF4F3	3C								
B→ / Aest_aes_new.128/a1/Aey_in B→ / Aest_aes_new.128/a1/Aey_in B→ / Aest_aes_new.128/a1/Aey_i0 (13 111 11 (100 14 F9 A8) (C2 SEC 25 89) (E1 3F 0C C8) (B6 63 0C A6)) F / Aest_aes_new.128/a1/Audu_0 H→ / Aest_aes_new.128/a1/Audu_0 H→ / Aest_aes_new.128/a1/Audu_fine_41/s_tox H→ / Aest_	Decryption														
B→       /test_ass_new.128/a1/key_00       (11 11 11       (10 0 14 F9 A8) (CS EF 25 93) (E1 3F 0C CS) (86 63 0 C A6))       Image: Comparison of the tass new.128/a1/dot_vaid_00         B→       /test_ass_new.128/a1/dot_vaid_00       (17 A D5 F (100 0 0 0) ((15 31 7 D E5) (0E 32 2C 72) (30 2 E 39 SF) (AF 09 07 94))       Image: Comparison of task new (128/a1/a01/ine_41/s_tow)         B→       /test_ass_new.128/a1/a01/ine_41/s_tow       (180 0 0 0) ((18 5 17 D E5) (0E 32 2C 72) (30 2 E 39 SF) (AF 09 07 94))       Image: Comparison of task new (128/a1/a01/ine_41/s_tow)         B→       /test_ass_new.128/a1/a01/ine_41/s_tow       (180 0 B C T (180 0 0 0) ((18 2 1 3 D2) (53 A1 42 1E) (38 C 3 F 2 84) (18 40 38 E7))       Image: Comparison of task new (180 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	⊞_ <td>00010203</td> <td>2B7E15</td> <td>1628AED2</td> <td>A6ABF715</td> <td>8809CF4F3</td> <td>3C</td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td>	00010203	2B7E15	1628AED2	A6ABF715	8809CF4F3	3C								
<ul> <li>Aret_aes_new_128/a1/dout_valid_00</li> <li>Round-1</li> <li>Round-1</li> <li>(7A D5 F (100 00 00)(1E3 31 70 B5) (1E3 22 C, 72) (30 2E 39 5F) (AF 09 07 54))</li> <li>Aret_aes_new_128/a1/A01/Ine_41/s_box</li> <li>(164 99 31 (26 27 26 F3) (19 FA DC 21) (28 D1 28 41) (55 5C 00 6F))</li> <li>(165 49 31 (26 27 26 F3) (19 FA DC 21) (28 D1 28 41) (55 5C 00 6F))</li> <li>(160 00 00 00) (00 bo)(1E8 40 F2 E1 (58 23 22 C, 72) (13 22 49 3F) (18 40 38 E7))</li> <li>(160 00 00 00) (00 bo)(1E8 40 F2 E1 (58 23 22 42 4) (18 40 38 E7))</li> <li>(160 00 00 00) (00 bo)(1E8 40 F2 E1 (58 23 22 4) (18 40 38 E7))</li> <li>(160 00 00 00) (00 bo)(1E8 40 F2 E1 (58 23 22 4) (18 40 38 E7))</li> <li>(161 47 73 B1 (100 00 00 00) (00 bo)(1E8 40 F2 E1 (58 20 2F))</li> <li>(164 74 73 B1 (100 00 00 00) (00 bo)(1E8 40 F2 E1 (58 20 2F))</li> <li>(164 74 73 B1 (100 00 00 00 00 00 00 00 00 00 00 00 00</li></ul>	<u>-</u> → /test_aes_new_128/a1/key_00	{{13 11 10	{{D014	F9 A8} {C9	EE 25 89	{E1 3F 0C	C8} {B6 6:	8 OC A6}}							
Bound-1       ((7A D5 F (00 00 00)((28 31 7D B5) (CE 32 C 72) (3D 2E 35 F) (AF 09 07 94))         B→ / hest_ase_new_128/a1/a01/ine_41/s_tow       ((7A D5 F (00 00 00 00))((EB 2E 13 D2) (59 A1 42 FE) (59 2C 38 41 (18 40 38 E7))         B→ / hest_ase_new_128/a1/a01/ine_41/m_cot       ((00 00 00 00) (00 00)((EB 40 F2 TE) (59 2C 38 41 (18 40 38 E7))         B→ / hest_ase_new_128/a1/a01/ine_41/m_cot       ((00 00 00 00) (00 00)((EB 40 F2 TE) (59 2C 38 41 (18 40 38 E7))         B→ / hest_ase_new_128/a1/a01/ine_41/m_cot       ((47 73 B) (00 00 00 00) (00 00)((EB 15 TC C5) (E 15 50 07 2) (FD A5 C 24) (F A0 38 21))         B→ / hest_ase_new_128/a1/a02/ine_41/m_cot       ((54 09 3) ((00 00 00 00) (00 00)((187 E4 6A 6A) (CE 83 42 55) (A9 44 4E AE) (0C F6 F5 59))         B→ / hest_ase_new_128/a1/a02/ine_41/m_cot       ((54 09 3) ((00 00 00 00) (00 00)((167 E4 6A 6A) (CE 83 42 55) (A9 44 4E AE) (0C F6 F5 59))         B→ / hest_ase_new_128/a1/a02/ine_41/m_cot       ((54 D3 9) ((00 00 00 00) (00 00 00) (00 00)((EA 35 5C F0) (04 45 33 2D) (65 56 5C 2D) (65 68 60 C5))         B→ / hest_ase_new_128/a1/a02/ine_41/m_cot       ((FD 28 8) ((00 00 00 00) (	🔶 /test_aes_new_128/a1/dout_valid_00	1													
B→       /test_aes_new_128/a1/ound_00       (7A D5F       {(100 00 00)({(E3 31 70 B5)}(CE 32 2C 72) (3D 2E 89 5F) (AF 03 07 94)}       Image: Control of Contro Control of Control of Control of Control of	Round-1														
g= 4       /test_aes_new_128/a1/key_01       ((54 99 3)       ((47 75 6 F3) (19 FA DC 21) (28 D1 28 41) (57 5 0 0 6E))	<u>-</u>	{{7A D5 F	{{00 00}}	00 )({E9	31 7D B5}	{CB 32 2C	72} {3D 2E	89 5F} {A	F 09 07 94)	}					
B→       /test_aes_new_128/a1/a01/line_41/s_tow       (BD B5 2       (BD 00 0 00) (D0 00) ((EE 2E 13 D2) (59 A1 42 1E) (8E C3 F2 84) (18 40 38 E7))         B→       /test_aes_new_128/a1/a01/line_41/m_col       (BD 05 1       (BD 00 0 00) (D0 00) ((EE 40 F2 TE) (59 Z5 38 84) (BE A1 13 E7) (18 C3 42 D2))       (BD 05 1         A→       /test_aes_new_128/a1/a01/line_41/m_col       (BD 05 00) (D0 00) ((EE 40 F2 TE) (59 Z5 38 84) (BE A0 38 21))       (BD 05 1         A→       /test_aes_new_128/a1/a01/line_41/m_col       (BD 05 00) (D0 00) ((EE 46 A6) F2 4C 176 D2 32 F))       (BD 05 100 00.0) ((BE 46 A6) F2 4C 176 D2 32 F))         B→       /test_aes_new_128/a1/a02/line_41/m_col       (BD 00 000) (D0 00) ((B1 46 46 A4) (CE 83 42 5) (A2 36 44 4E AE) (0C FE F5 59)       (BD 05 100 000 000) (D0 00 00 00) (D0		{{54 99 33	{{AC 77	66 F3} {19	FA DC 21	{28 D1 29	41} (57 50	: 00 6E}}							
B→       /test_aes_new_128/a1/a01/line_41/s_row       (BD 6E 7       ((00 00 00 00) (00 00)((EE 40 F2 1E) (59 2E 38 64) (8B A1 13 E7) (18 C3 42 D2))         B→       /test_aes_new_128/a1/dout_vaid_01       ((47 78 8)       ((00 00 00 00) (00 00)((EE 40 F2 1E) (59 2E 38 64) (8B A1 13 E7) (18 C3 42 D2))         B→       /test_aes_new_128/a1/dout_vaid_01       ((47 78 8)       ((10 00 00 00) (00 00)((8E 15 1C C5) (E1 55 00 72) (FD A2 C2 48) (F A0 38 21))       (10 00 00 00) (00 00)((8E 16 1C C5) (E1 55 00 72) (FD A2 C2 48) (F A0 38 21))         B→       /test_aes_new_128/a1/key_02       ((47 43 8)       ((EA D2 73 21) (EE 50 BA D2) (31 2E F5 60) (7F 8) 29 2F).)       (10 00 00 00) (00 00 00)		{(BD B5 2	{{00 00}}	00 00} {00	00 <u>)</u> {{EB	2E 13 D2}	{59 A1 42	1E} {8B C3	F2 84} {1E	40 38 E 7	8				
g= 4       /test_aes_new_128/a1/a01/line_41/m_col       {(47 73 B       {(00 00 00 00)		{{BD 6E 7	{{00 00}}	00 00} {00	00 ){{EB	40 F2 1E}	{59 2E 38	84} {8B A1	13 E7} {1E	C3 42 D2	8				
		{{47 73 B	{{00 00}	00 00} {00	00 )({{8B	15 1C C5}	{E1 55 0D	72} {FD A9	C2 48} {F	A0 38 21	8				
Bound-2       Image: provide	/test_aes_new_128/a1/dout_valid_01	1													
P= ^ / Ast_ass_new_128/a1/round_01       ((54 D 9)       ((00 00 0 0 00) (00 00)(((37 E4 64 64) F2 4C E7 8C) (4D 90 44 08) (37 EC C3 95))         P= ^ / Ast_ass_new_128/a1/a02/ine_41/m_col       ((47 43 8)       ((16 0 0 7 E 8)       ((17 8 6) 23 27))       ((17 8 6) 23 27)         P= ^ / Ast_ass_new_128/a1/a02/ine_41/s_tow       ((F0 E3 8)       ((00 0 0 0 0 0))       ((00 0 0 0 0 0))       ((17 8 6) 23 27))       ((23 8 8) 0 (25 8) 8 60 C5))         P= ^ / Ast_ass_new_128/a1/a02/ine_41/s_tow       ((F0 E5 9)       ((00 0 0 0 0 0))       ((00 0 0 0 0 0))       ((24 43 32 C))       ((25 9 8 40))       (25 9 8 50 C5))       ((24 5 33 2C))       ((25 9 8 60 C5))       ((20 0 0 0 0 0))       ((17 8 6) 23 22))       ((25 9 8 60 C5))       ((26 0 0 0 0 0))       ((17 8 6) 23 22))       ((25 9 8 60 C5))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26 0 0 0 0 0 0))       ((26															
B→       /test_aes_new_128/a1/key_02       (47 43 8       (47 48 8 3 2.5)       (4		{{54 D9 9	{{00 00}}	00 00} {00	00 )({87	6E 46 A6}	(F2 4C E7	BC} {4D 90	4A D8} {9	7 EC C3 95	}}				
		{{47 43 8	{{EA D2	73 21} {B	8D BA D	2) {31 2B F	5 60} {7F 8	D 29 2F}}							
Image: A / Ast_ass_new_128/a1/a02/line_41/s_tow       ((FD E3 8 (000 00 00 00) (00 00 00		{{2D 7E 8	{{00.00	00 00} {00	{00 00 00}	100	46 46 A4}	CB 83 42 5	55} {A9 44	4E AE} {0C	F6 F5 69}				
B+       /test_aes_new_128/a1/a02/line_41/s_box       ((FD E5 3 (00 00 00 00) (00 00 00) (00 00 00) (00 00 00 00) (00 00		{{FD E3 B	{{00.00	00 00} {00	00 00 00}	00){{EA	83 5C FO}	04 45 33 2	2D} {65 5D	98 AD} {85	5 96 BO C5	}			
		{{FD E5 9	{{00.00	00 00} {00	00 00 00}	00. ){{EA	45 98 C5}	04 5D BO	F0} {65 96	5C 2D} {85	83 33 AD	}			
■ Acund-10       ■ <td< td=""><td>/test aes new 128/a1/dout valid 02</td><td>1</td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td></td<>	/test aes new 128/a1/dout valid 02	1													
Image: Answing 128/a1/alast/line_30/s_box       ((00 50 All (00 00 00 00) (00 b0 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00) (00 00 00		{{63 53 E	{{00.00}	00 00} {00	00 00 00}	(00 00 00 C	00 00 101	0 00}}					X{D4P	(F 5D 30) {	E0 B4
	TH- /test aes new 128/a1/alast/line 30/s box	{{00 50 A	{{00.00	00.00} (00	00.00.00}		103 {00 00 (	0.00}}						¥{19 F	4 8D 0
B→       /test_ass_new_128/a1/key_lst       ((00 01 0)       ((28 7c [15 16] (28 Ac D2 A6] (AB F7 15 88) (03 C [24 73 C))       ((00 10 10)       ((28 7c [15 16] (28 Ac D2 A6] (AB F7 15 88) (03 C [24 73 C))         B→       /test_ass_new_128/a1/dout_lst       ((00 11 2)       ((00 00 00 00) (00 00 00 00) (00 00 00 00))       ((00 10 12)       ((00 11 2)       ((00 11 2)       ((00 10 00 00 00) (00 00 00))       ((00 10 00 00 00))       ((01 12)       ((01 12)       ((00 10 00 00 00))       ((01 12)		{{00.10.20	{{00.00	00.003.(00	00.00.00}		00 00 00 00	0.00}}						Y{(193	DE3B
B         //test_aes_new_128/a1/dout_ist         ((00 11 27) (00 00 00 00)         (00 00 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10 00)         ((00 10	The states new 128/a1/key lst	800.01.02	{{2B 7E	15 16} {28	ΔE D2 Δ6	(AB F7 15	88) (09 01	E 4E 3033							0 20 0
	The states new 128/a1/dout 1st	800 11 22	{{00.00	nn nn} (nn	00.00.003		n) {nn nn r	10 0033						Y832 4	3 E6 A
Final PlainText     ☐	/test_aes_new_128/a1/dout_valid	1	((00 00	00 00) (00	00 00 00)		0,00000	,0 00,)							0104
B→         //test_aes_new_128/a1/data_out         00112235         000000000000000000000000000000000000	Final PlainText														
	The states new 128/a1/data out	00112233	000000											13243E	64888
140 100 100 200 200 240	E A Host cost went i con any data four	00172200	0000000		000000000000000000000000000000000000000										
140 160 180 200 220 240				1.	40	11	50	1	80	21		22	0	24(	3

Figure 10. Case 4: Internal Operation of decoding phase of proposed architecture

CASE-3: This case presents the detail analysis of the internal operation performed by the decoding system of the AES. Each of the transformation with input and corresponding output of each round could be seen in the figure 10. The decoding process incorporates a series of transformations such as inv\_substitution, inv\_shifting and inv\_mixing of columns to obscure the information.



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	iiuc. 5.00											Jaci I'd	. 101.1	
	/test_aes_new_128/a0/ce	1												t
	/test_aes_new_128/a0/clk	0												L
	/test_aes_new_128/a0/reset	0												
— E	ncryption													
	/test_aes_new_128/a0/enc_dec	1												F
	/test_aes_new_128/a0/din_valid	1												F
<b>H</b> -4	/test_aes_new_128/a0/data_in	001122	(3243F	648885430	8D313198A2	E0370734								F
	/test_aes_new_128/a0/k_en	0												
<b>H</b> -4	/test_aes_new_128/a0/key_in	000102	000)(2B	7E151628AE	D2A6ABF71	58809CF4F3	C							F
<b>H</b> -4	/test_aes_new_128/a0/key_00	{{00 01	{{0 }{{2	B 7E 15 16}	{28 AE D2 A	6} {AB F7 15	88} {09 CF ·	4F 3C}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_00	{{00 10	{{00 00 0}	)({{19 3D I	3 BE} {A0 F	4 E 2 2B} {9A	C6 8D 2A}	E9 F8 48 08	}}					F
<b>H</b> -4	/test_aes_new_128/a0/key_01	{{D6 A4	{{6 }{{{	0 FA FE 17}	{88 54 2C B	1} {23 A3 39	39} {2A 6C 7	6 05}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_01	{{89 D8	{{00 00 0	0 00} { )({{	4 9C 7F F2}	{68 9F 35 2E	} {6B 5B EA	43} {02 6A 5	0 49}}					F
<b>H</b> -4	/test_aes_new_128/a0/key_02	{{B6 92	{{9 )({{F	2 C2 95 F2}	{7A 96 B9 43	8} {59 35 80 1	7A} {73 59 FI	6 7F}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_02	{{49 15	{{00 00 0	0 00} {00 00	){{{AA 8F !	5F 03} {61 D(	) E3 EF} {82	D2 4A D2} {	68 32 46 <del>3</del> A	}}				F
<b>H</b> -4	/test_aes_new_128/a0/key_03	{{B6 FF	{{9 }{{3	D 80 47 7D)	{47 16 FE 3	E} {1E 23 7E	44} {6D 7A	88 3B}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_03	{{FA 63	{{00 00 0	0 00} {00 00	00 00} )({{4	8 6C 4E EE}	{67 1D 9D 0	D} {4D E3 B	1 38} {D6 5F	58 E 7}}				F
<b>H</b> -4	/test_aes_new_128/a0/key_04	{{47 F7	{{E )({{E	F 44 A5 41}	{A8 52 5B 7	F} {B6 71 25	3B} {DB 0B /	AD 00}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_04	{{24 72	{{00 00 0	0 00} {00 00	00 00} {00 00	)({{E0 92 7	F E8} {C8 63	3 63 C0} {D9	B1 35 50} {8	5 B8 BE 01}	}			F
<b>H</b> -4	/test_aes_new_128/a0/key_05	{{3C AA	{{7 }{{[	4 D1 C6 F8)	{7C 83 9D 8	7} {CA F2 B8	BC} {11 F9	15 BC}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_05	{{C8 16	{{00 00 0}	0 00} {00 00	00 00} {00 0	D OO O )({{F	1 00 6F 55}	[C1 92 4C EF	F} {7C C8 8B	32} {5D B5 [	05 OC}}			
<b>H</b> -4	/test_aes_new_128/a0/key_06	{{5E 39	{{E }{{E	D 88 A3 7A)	{11 OB 3E F	D} {DB F9 86	641} (CA 00	93 FD}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_06	{{C6 2F	{{00 00 0}	0 00} {00 00	00 00} {00 0	00 00 {00 00 {	)({{26 OE 2	E 17} (3D 4	1 B7 7D} {E8	64 72 A9} {F	D D2 8B 25	}}		
<b>—</b>	/test_aes_new_128/a0/key_07	{{14 F9	{{2 }{{4	E 54 F7 0E}	{5F 5F C9 F3	8} {84 A6 4F	82} {4E A6 E	)C 4F}}						F
<b>H</b> -4	/test_aes_new_128/a0/round_07	{{D1 87	{{00 00 0}	0 00} {00 00	00 00} {00 0	00 00 {00 00 {	00 00 )({{5	A 41 42 B1}	{19 49 DC 1	F} {A3 E0 19	65} {7A 8C (	)4 OC}}		
<b>H</b> -4	/test_aes_new_128/a0/key_08	{{47 43	{{0 }({{E	A D2 73 21)	(B5 8D BA I	2} {31 2B F	5 60} {7F 8D	29 2F}}						
<b>H</b> -4	/test_aes_new_128/a0/round_08	{{FD E3	{{00 00 0}	0 00} {00 00	00 00} {00 0	00 00 {00 00 {	00 00 00}}	){{EA 83 5	5C F0} {04 4!	33 2D} {65	5D 98 AD} {	35 96 BO C5}	}	
<b>H</b> -4	/test_aes_new_128/a0/key_09	{{54 99	{{B }{{{A	C 77 66 F3}	{19 FA DC 2	1} {28 D1 29	41} {57 5C 0	0 6E}}						
<b>H</b> -4	/test_aes_new_128/a0/round_09	{{BD 6E	{{00 00 0}	0 00} {00 00	00 00} {00 0	00 00 {00 00 {	00 00 00}}	){{{E	B 40 F2 1E}	{59 2E 38 84	\$} {8B A1 13	E7} {1B C3 4	2 D2}}	
<b>H</b> -4	/test_aes_new_128/a0/key_lst	{{13 11	{{B }{{{C	0 14 F9 A8}	{C9 EE 25 8	9} {E1 3F 0C	C8} {B6 63	OC A6}}						
<b>—</b>	/test_aes_new_128/a0/dout_lst	{{69 C4	{{00 00 0	0 00} {00 00	00 00} {00 0	0 00 00} {00	00 00 00}}		){{39 25 8	4 1D} {02 D0	C 09 FB} {DC	11 85 97} {1	9 6A 0B 32}	F
- 4	/test_aes_new_128/a0/dout_valid	1												
<b>—</b>	/test_aes_new_128/a0/data_out	69C4E0	00000000	000000000000000000000000000000000000000	000000000000000000000000000000000000000	00000			(3925841	DO2DCO9FB	DC11859719	6A0B32		F
			0.00	i la co			1.1.1							t
					50		1(	00		150		20	<u> </u>	

Figure 11. Encryption process of proposed architecture

Case-4: This case presents the detail analysis of the entire encryption and decryption process of the AES with outputs at each stage is presented in figures 11 and figure 12 respectively with input data from case-1.

Decyption														
🔶 /test_aes_new_128/a1/enc_dec	0													
	1													
⊞	69C4E0	(3925)	341D02DC0	9FBDC1185	97196A0B33	2								
⊞	000102	2B7E15	1628AED24	6ABF71588	D9CF4F3C									
⊞	{{13 11	{{D0 14	F9 A8} {C9	EE 25 89} {	1 3F OC C8	} {B6 63 OC.	A6}}							
· ⊕	{{7A D5	{{00 00}}	00 )({{E9 :	31 7D B5} {(	B 32 2C 72	{3D 2E 89	5F} {AF 09 0	7 94}}						
⊞ /test_aes_new_128/a1/key_01	{{54 99	{{AC 77	66 F3} {19 I	A DC 21} {:	28 D1 29 41	{57 5C 00 (	5E}}							
· ⊞	{{54 D9	{{00 00}}	00 00} {00 00	0 0 ){{{87 6	SE 46 A6} {F	2 4C E7 8C)	{4D 90 4A [	97 EC 0	3 95}}					
⊞–♦ /test_aes_new_128/a1/key_02	{{47 43	{{EA D2	73 21} {B5	8D BA D2}-	(31 2B F5 60	)} {7F 8D 29	2F}}							
· ⊕	{{3E 1C	{{00 00}}	00 00} {00 00	0 00 00} {0(	00){{BE	3B D4 FE} {	D4 E1 F2 C8	} {0A 64 2C	CO} {DA 83	86 4D}}				
⊞–♦ /test_aes_new_128/a1/key_03	{{14 F9	{{4E 54	F7 0E} {5F !	5F C9 F3} {8	4 A6 4F B2}	{4E A6 DC -	4F}}							
· ⊞	{{B4 58	{{00 00}}	00 00} {00 00	0 00 00} {0(	00 00 00}	(00 )({{F7 8	3 40 3F} {2	7 43 3D FO}	(98 B5 31 Fi	F} {54 AB AS	9 D 3}}			
⊞–♦ /test_aes_new_128/a1/key_04	{{5E 39	{{6D 88	A3 7A} {11	0B 3E FD}{	DB F9 86 41	} {CA 00 93	FD}}							
· ⊞	{{E8 DA	{{00 00}}	00 00} {00 00	0 00 00} {0(	00 00 00}	00 00 00 00	}} )({{A1 4	4F 3D FE} {7	8 E8 03 FC)	{10 D5 A8	DF} {4C 63 2	9 23}}		
⊞–♦ /test_aes_new_128/a1/key_05	{{3C AA	{{D4D1	C6 F8} {7C	83 9D 87} {	CA F2 B8 B0	;} {11 F9 15	BC}}							
⊞_ <td>{{36 33</td> <td>{{00 00}}</td> <td>00 00} {00 0</td> <td>0 00 00} {0(</td> <td>00 00 00 00</td> <td>00 00 00 00</td> <td>}}</td> <td>)({{E1 </td> <td>FB 96 7C} {E</td> <td>8 C8 AE 9B</td> <td>} {35 6C D2</td> <td>BA} {97 4F F</td> <td>B 53}}</td> <td></td>	{{36 33	{{00 00}}	00 00} {00 0	0 00 00} {0(	00 00 00 00	00 00 00 00	}}	)({{E1	FB 96 7C} {E	8 C8 AE 9B	} {35 6C D2	BA} {97 4F F	B 53}}	
⊞	{{47 F7	{{EF 44	A5 41} {A8	52 5B 7F} {E	6 71 25 3B)	{DB 0B AD	00}}							
⊞_ <td>{{2D 6D</td> <td>{{00 00}}</td> <td>00 00} {00 0</td> <td>0 00 00} {0(</td> <td>00 00 00 00</td> <td>00 00 00 00</td> <td>}}</td> <td></td> <td>)({{527</td> <td>4 C8 94} {8</td> <td>5 11 6A 28}</td> <td>{E3 CF 2F D</td> <td>7} {F6 50 5</td> <td>07}}</td>	{{2D 6D	{{00 00}}	00 00} {00 0	0 00 00} {0(	00 00 00 00	00 00 00 00	}}		)({{527	4 C8 94} {8	5 11 6A 28}	{E3 CF 2F D	7} {F6 50 5	07}}
⊞	{{B6 FF	{{3D 80	47 7D} {47	16 FE 3E} {	IE 23 7E 44	{6D 7A 88	3B}}							
⊞_ <td>{{3B D9</td> <td>{{00 00}}</td> <td>00 00} {00 0</td> <td>0 00 00} {0(</td> <td>00 00 00 00</td> <td>00 00 00 00</td> <td>}}</td> <td></td> <td></td> <td>){{AC</td> <td>C1 D6 B8} {</td> <td>EF B5 5A 7B</td> <td>} {13 23 CF</td> <td>DF} {45 7</td>	{{3B D9	{{00 00}}	00 00} {00 0	0 00 00} {0(	00 00 00 00	00 00 00 00	}}			){{AC	C1 D6 B8} {	EF B5 5A 7B	} {13 23 CF	DF} {45 7
⊞	{{B6 92	{{F2 C2	95 F2} {7A !	6 B9 43} {5	9 35 80 7A}	{73 59 F6 7	F}}							
· ⊕	{{A7 BE	{{00 00}}	00 00} {00 0	0 00 00} {0(	00 00 00 00	00 00 00 00	}}				){{{49 [	)B 87 3B}{4	5 39 53 89}	{7F 02 D ;
⊞–✦ /test_aes_new_128/a1/key_09	{{D6 A4	{{A0 FA	FE 17} {88	54 2C B1} {2	3 A3 39 39)	{2A 6C 76 0	)5}}							
· ⊕	{{63 53	{{00 00}}	00 00} {00 0	0 00 00} {0(	00 00 00 00	00 00 00 00	}}					){{D4	BF 5D 30} {E	0 B4 52 )
⊞	{{00 01	{{2B 7E	15 16} {287	AE D2 A6} {	AB F7 15 88	} {09 CF 4F :	3C}}							
⊕	{{00 11	{{00 00}}	00 00} {00 0	0 00 00} {00	00 00 00 00	00 00 00 00	}}						){{{32 4	3 F6 A8}
/test_aes_new_128/a1/dout_valid	1													
· ⊞–- /test_aes_new_128/a1/data_out	001122	000000	0000000000	0000000000	000000								(3243	6A8885
			14	40	11	50	18	B0	20	)0	2.	20	24	40

Figure 12. Decryption process of proposed architecture

# CONCLUSION

In this thesis, we presented a VLSI based AES (Advanced Encryption Standard) encryption that effectively addressed espionage and fraudulent cybercrime based cyber attacks. It is most commonly used symmetric block cipher algorithm that transform information into obscure data based on key-defined transformation set in an effective and robust manner. In addition, it is lossless operation (i.e. original information and decrypted information are similar with no

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distortion) with size of input and output being the same and could be extended to a wide range of applications. We limit our focus to128-bit AES encoding and decoding operations over VHDL coded transformation that requires key for successful completion of the operation

With rise in several methods proposed for implementation of data security it has become more vital for a feasibility study of any hardware design is essential i.e. to test encryption and decryption process of the proposed 128-bit AES algorithm. Successful implementation and feasibility analysis of the AES encoding and decoding algorithm are carried out using ModelSim. In addition, the simulations results prove that the proposed implement offers promising results in speed and delay in comparison with other schematics. The analysis of the each of the transformation was also performed based coding on FPGA using Xilinx ISE tool.

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